

Evolutionary Computation

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Course organization

- Part 1: lectures \Rightarrow Dirk Thierens + 2 guest lectures Peter Bosman
- Part 2: seminar \Rightarrow papers & presentation (student groups)
- Part 3: practical assignment \Rightarrow report (groups of 1 or 2 students)
 - ▶ Specific discrete benchmark functions
 - ▶ Graph Bipartitioning

Course grading

- 1 Written exam = 40%
- 2 Paper presentation = 20%
- 3 Practical assignment report = 40%

Evolutionary Computation

**= Population-based, stochastic search algorithms
inspired by mechanisms of natural evolution**

- EC part of Computational Intelligence
- Evolution viewed as search algorithm
- Natural evolution only used as metaphor for designing computational problem solving systems
- No modelling of natural evolution (\neq evolutionary biology)

Evolutionary algorithm

```

1 P(0) ← Generate-Random-Population()
2 P(0) ← Evaluate-Population(P(0))
3 while Not-Terminated? do
  1 Ps(t) ← Select-Mates(P(t))
  2 Po(t) ← Generate-Offspring(Ps(t))
  3 Po(t) ← Evaluate-Population(Po(t))
  4 P(t+1) ← Select-Fittest(Po(t) ∪ P(t))
  5 t ← t + 1
4 return P(t)

```

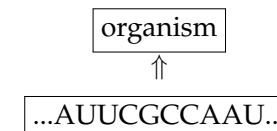
Darwinian process characteristics ⇒ GA

- 1 Structures
⇒ e.g. **binary strings, real-valued vectors, programs, ...**
- 2 Structures are copied
⇒ **selection algorithm**
- 3 Copies partially vary from the original
⇒ **mutation & crossover operators**
- 4 Structures are competing for a limited resource
⇒ **selecting fixed sized parent pool**
- 5 Reproductive success depends on environment
⇒ **user defined fitness function**

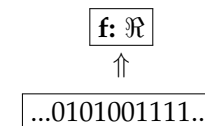
Key concepts of a Darwinian system

- 1 Information Structures
- 2 Copies
- 3 Variation
- 4 Inheritance
- 5 Competition

• Neo-Darwinism



• Genetic Algorithm



- * user: string representation and function f
- * GA: string manipulation
 - ▶ **selection:** copy better strings
 - ▶ **variation:** generate new strings

- **selection:** copy better strings
 - ▶ tournament selection
 - ▶ truncation selection
 - ▶ proportionate selection
- **variation:** generate new strings

1 crossover

$$\text{2-point crossover: } \begin{cases} 1111111111 \\ 0000000000 \end{cases} \Rightarrow \begin{cases} 1111000011 \\ 0000111100 \end{cases}$$

$$\text{uniform crossover: } \begin{cases} 1111111111 \\ 0000000000 \end{cases} \Rightarrow \begin{cases} 1001110101 \\ 0110001010 \end{cases}$$

2 mutation

$$\{1111111111\} \Rightarrow \{1111111011\}$$

- **Generation 1:**
tournament selection, 1-point crossover, mutation

Parents	Fitness	Offspring	Fitness
100!10	324	10 <u>1</u> 00	400
101!00	400	1011 <u>1</u>	529
01!000	64	00 <u>0</u> 10	4
10!010	324	100 <u>1</u> 0	324
0110!0	144	<u>1</u> 1100	784
1010!0	400	10 <u>0</u> 00	256

Parent population mean fitness $\bar{f}(1) = 383$

Toy example

$$x \in [0, 31] : f(x) = x^2$$

binary integer representation: $x_i \in \{0, 1\}$

$$x = x_1 * 2^4 + x_2 * 2^3 + x_3 * 2^2 + x_4 * 2^1 + x_5 * 2^0$$

- Initial Random Population:

$$10010 : 18^2 = 324$$

$$01100 : 12^2 = 144$$

$$01001 : 9^2 = 81$$

$$10100 : 20^2 = 400$$

$$01000 : 8^2 = 64$$

$$00111 : 7^2 = 49$$

population mean fitness $\bar{f}(0) = 177$

- **Generation 3:**

Parents	Fitness	Offspring	Fitness
1!1111	961	111 <u>1</u> 0	900
1!1100	784	110 <u>1</u> 1	729
110!00	576	11 <u>1</u> 10	900
111!10	900	1110 <u>1</u>	841
1101!1	729	11 <u>1</u> 11	961
1100!1	625	0 <u>1</u> 001	81

Parent population mean fitness $\bar{f}(0) = 762$

Schema = similarity subset

eg. : $11\#\#0 = \{11000, 11010, 11100, 11110\}$

gen.	1####	0####	####1	####0
0	2	4	2	4
1	5	1	1	5
2	6	0	2	4
3	6	0	3	3
4	6	0	3	3
5	5	1	4	2

Schema

- definitions:
 - ▶ $o(h)$: schema order $o(11\#\#0) = 3$
 - ▶ $\delta(h)$: schema defining length $\delta(11\#\#0) = 4$
 - ▶ $m(h,t)$: number of schema h instances at generation t
 - ▶ $f(h,t) = \frac{\sum_{i \in P} f_i}{m(h,t)}$: schema fitness is average fitness of individual members
- key issue: **changing number of schemata members in population**
- fit schemata increase in proportion
- mutation and recombination destructive operators !

Schema growth by selection

- Reproduction ratio $\phi(h,t)$

$$\phi(h,t) = \frac{m(h,t^s)}{m(h,t)}$$

- **proportionate selection**

- ▶ probability individual i selected: $\frac{f_i}{\bar{f}(t)}$
- ▶ Expected number of copies that are member of schema h after selection:

$$m(h,t^s) = m(h,t)\phi(h,t) = m(h,t)\frac{f(h,t)}{\bar{f}(t)}$$

- **tournament selection**

- ▶ tournament size s : $0 \leq \phi(h,t) \leq s$

Schema disruption by mutation

- probability bit flipped: p_m
- schema h survives iff all the bit values are *not* mutated

$$p_{survival} = (1 - p_m)^{o(h)}$$

- for small values $p_m \ll 1$

$$(1 - p_m)^{o(h)} \approx 1 - o(h) \cdot p_m$$

- disruption factor $\epsilon(h,t)$ by mutation:

$$\epsilon(h,t) = o(h) \cdot p_m$$

Schema disruption by recombination

- probability crossover applied p_c
- **1-point crossover**
 - ▶ schema h survives iff cutpoint *not* within defining length δ :

$$p_{\text{survival}} = 1 - \frac{\delta(h, t)}{l - 1}$$

- **uniform crossover** (bit swap probability: p_x)
 - ▶ schema h survives iff none or all bits swapped together

$$p_{\text{survival}} = p_x^{o(h)} + (1 - p_x)^{o(h)}$$

- disruption factor $\epsilon(h, t)$ by recombination:

$$\epsilon(h, t) = p_c \cdot (1 - p_{\text{survival}})$$

Schema Theorem

- Selection, mutation, and recombination combined:

$$m(h, t + 1) \geq m(h, t) \phi(h, t) [1 - \epsilon(h, t)]$$

- net growth factor: $\gamma(h, t) = \frac{m(h, t+1)}{m(h, t)}$

$$\gamma(h, t) \geq \phi(h, t) [1 - \epsilon(h, t)]$$

schemata with $\gamma(h, t) > 1$ increase in proportion
 schemata with $\gamma(h, t) < 1$ decrease in proportion

Schema Theorem cont'd

- low order, high performance schemata receive exponentially (geometrically) increasing trials → **building blocks**
- according to the k-armed bandit analogy this strategy is near optimal (Holland, 1975)
- happens in an implicit parallel way
 - only the short, low-order schemata are processed reliably
- enough samples present for statistically reliable information
- enough samples survive the disruption of genetic operators

Building Blocks

Building block hypothesis

= building blocks can be juxtaposed to form near optimal solutions

Consequences

- 1 schema sampling is a statistical decision process:
variance considerations
- 2 building blocks must be juxtaposed before convergence:
mixing analysis
- 3 low order schemata might give misleading information:
deceptive problems

Permutation problems

- **Goal**
Design suitable representations and genetic operators for permutation or sequencing problems
- **Examples**
 - ▶ scheduling
 - ▶ vehicle routing
 - ▶ queueing
 - ▶ ...

Permutation problems

- travelling salesman
- non-binary strings
 - ▶ $p1 = 1\ 2\ 3\ 4\ 5\ 6\ 7\ 8$
 - ▶ $p2 = 4\ 6\ 2\ 1\ 7\ 8\ 5\ 3$
 - ▶ simple crossover \Rightarrow illegal tours
 - ▶ $c1 = \underline{1}\ 2\ 3\ | \underline{1}\ 7\ 8\ 5\ 3$
 - ▶ $c2 = \underline{4}\ \underline{6}\ 2\ | \underline{4}\ 5\ \underline{6}\ 7\ 8$
- alternative search space representation
- alternative genetic operators

Insert mutation

randomly select one element from the sequence and insert it at some other random position in the sequence

$$\begin{array}{cccccccc} A & B & \underline{C} & D & E & F & G & H \\ & & \Downarrow & & & & & \\ A & B & D & E & F & \underline{C} & G & H \end{array}$$

Swap mutation

randomly select two elements from the sequence and swap their position

$$\begin{array}{cccccccc} A & B & \underline{C} & D & E & F & \underline{G} & H \\ & & \Downarrow & & & & \Downarrow & \\ A & B & \underline{G} & D & E & F & \underline{C} & H \end{array}$$

Scramble mutation

randomly select a subsequence and scramble all elements in this subsequence

$$\begin{array}{cccccc|cccc} A & B & & C & D & E & F & & G & H \\ & & & & & \downarrow & & & & \\ A & B & & D & F & E & C & & G & H \end{array}$$

very destructive !

→ efficiency is problem dependent !

Mutation operator: 2-opt

⇒ randomly select two points along the sequence and invert one of the subsequences

$$\begin{array}{cccccc|cccc} A & B & & C & D & E & F & & G & H \\ & & & & & \downarrow & & & & \\ A & B & & F & E & D & C & & G & H \end{array}$$

2-opt can be applied to $\frac{n(n-1)}{2}$ pairs of edges: if none of these gives an improvement a local optimum has been reached.

Mutation operators

- TSP: *adjacency* of elements in permutation is important
→ 2-opt only minimal change
- scheduling: *relative ordering* of elements in permutation is important
→ 2-opt large change
e.g.: priority queue: line of people waiting for supply of tickets for different seats on different trains

mutation principle: “small” moves in search space should be more likely than “large” moves

Recombination operators

- ‘standard’ crossover operators generate infeasible sequences

$$\begin{array}{cccccc|cccc} A & B & C & D & E & & F & G & H \\ b & f & d & h & g & & e & a & c \\ & & & & \downarrow & & & & \\ A & B & C & D & E & & e & a & c \\ b & f & d & h & g & & F & G & H \end{array}$$

- different aspects
 - ▶ adjacency
 - ▶ relative order
 - ▶ absolute order

⇒ whole set of permutation crossover operators proposed !

Order crossover

p1: A B | C D E F | G H I
 p2: h d | a e i c | f b g
 ↓
 ch: a i C D E F b g h

- 1 randomly select two crosspoints
- 2 copy subsequence between crosspoints from p1
- 3 starting at 2nd crosspoint: fill in missing elements retaining relative order from p2

Partially mapped crossover

p1: A B | C D E F | G H I
 p2: h d | a e i c | f b g
 ↓
 ch: h i C D E F a b g

- 1 randomly select two crosspoints
- 2 copy p2 to child
- 3 copy elements between crosspoints from p1 to child while placing the replaced element from p2 at the location where the replacer is positioned

Position crossover

p1: A B C D E F G H I
 p2: h d a e i c f b g
 * * * *
 ↓
 ch: A h C d E F b g I

- 1 randomly select k positions
- 2 copy unmarked elements from p1 to child
- 3 scan p2 from left to right and fill in missing elements

Maximal preservative crossover

p1: A B | C D E F | G H I
 p2: h d | a e i c | f b g
 ↓
 ch: i a C D E F b g h

- 1 randomly select two crosspoints
- 2 copy subsequence between crosspoints from p1
- 3 add successively an adjacent element from p2 starting at last element in child
- 4 if already placed: take adjacent element from p1

Cycle crossover

p1: A B C D E F G H I
 p2: f c d a e b h i g
 cy: 1 1 1 1 2 1 3 3 3
 ↓
 ch: A B C D E F h i g

- 1 mark cycles
- 2 cross full cycles

⇒ emphasizes absolute position above adjacency or relative order

edge recombination

parent tours [ABCDEF] & [BDCAEF]

edge map:

city	edges
A	B F C E
B	A C D F
C	B D A
D	C E B
E	D F A
F	A E B

edge recombination algorithm:

- 1 choose initial city from one parent
- 2 remove current city from edge map
- 3 if current city has remaining edges
goto step 4
else
goto step 5
- 4 choose current city edge with fewest remaining edges
- 5 if still remaining cities, choose one with fewest remaining cities

- 1 random choice ⇒ B
- 2 next candidates: A C D F
choose from C D F (same edge number) ⇒ C
- 3 next candidates: A D
(edgelist D < edgelist A) ⇒ D
- 4 next candidate: E ⇒ E
- 5 next candidates: A F
tie breaking ⇒ A
- 6 next candidate: F ⇒ F

resulting tour: [BCDEAF]

Fitness correlation coefficients

- genetic operators should preserve useful fitness characteristics between parents and offspring
- calculate the fitness correlation coefficient to quantify this
- k-ary operator: generate n sets of k parents
- apply operator to each set to create children
- compute fitness of all individuals
- $\{f(p_{g1}), f(p_{g2}), \dots, f(p_{gn})\}$
- $\{f(c_{g1}), f(c_{g2}), \dots, f(c_{gn})\}$

Fitness correlation coefficients

- F_p : mean fitness of the parents
 F_c : mean fitness of the children
 $\sigma(F_p)$ = standard deviation of fitness parents
 $\sigma(F_c)$ = standard deviation of fitness children
 $cov(F_p, F_c) = \sum_{i=1}^n \frac{(f(p_{gi}) - F_p)(f(c_{gi}) - F_c)}{n}$
covariance between fitness parents and fitness children
- operator fitness correlation coefficient ρ_{op} :

$$\rho_{op} = \frac{cov(F_p, F_c)}{\sigma(F_p)\sigma(F_c)}$$

Traveling Salesman problem: mutation operators

- various mutation operators applicable
 - 2opt mutation (*2OPT*)
 - swap mutation (*SWAP*)
 - insert mutation (*INS*)

performance: *2OPT* > *INS* > *SWAP*

- mutation fitness correlation coefficients ρ_{mutate} :

ρ_{2OPT}	0.86
ρ_{INS}	0.80
ρ_{SWAP}	0.77

Traveling Salesman problem: crossover operators

- various crossover operators in applicable
 - cycle crossover (*CX*)
 - partially matched crossover (*PMX*)
 - order crossover (*OX*)
 - edge crossover (*EX*)

performance: *EX* > *OX* > *PMX* > *CX*

- crossover correlation coefficients ρ_{cross} :

ρ_{EX}	0.90
ρ_{OX}	0.72
ρ_{PMX}	0.61
ρ_{CX}	0.57