Soft Typing for JavaScript

Liewe Thomas van Binsbergen João Paulo Pizani Flor Pepijn Kokke

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Parsing JavaScript

Using Language. JavaScript¹

We started out using the Language. JavaScript package, after it was approved with the remark that with all the clutter of lexical information the AST made the parser look more like a lexer than an actual parser.

Sadly, this turned out to be closer to the truth than we originally imagined. After the creation of a module that removed all the lexical clutter from the ${\rm AST^2}$, the tree we ended up with contained many structures such as the following.

These structures would require a parser of their own to interpret into sensible expressions—an issue that will only show once you try to use the JavaScript AST. However, since the objective of Language.JavaScript was to create a valid JavaScript parser in Haskell, this was probably never noticed.

For this reason we decided to move to a parser that, while it hasn't been worked on in quite some time, was actually used to write a JavaScript interpreter.

Using the HJS Parser³

The HJS parser—as opposed to Language.JavaScript—does group expressions according to their semantics. This means, for instance, that the assignment

¹See http://hackage.haskell.org/package/language-javascript-0.5.7.

²This module is no longer part of the package, but may be found in the histories of our git commit logs.

³See http://hackage.haskell.org/package/hjs-0.2.1.

operators are no longer represented as generic binary operators, but as a different class of statements.

This made the interpretation of the HJS AST much, much simpler, and in mere hours we were able to make more progress that we had made in the weeks before.

However, the semantic separation that HJS imposes causes the AST to become much more complex, as it makes distinctions such as effectful and non-effectful, prefix or postfix, etc...

We were able to convert a large number of these expressions to our analysis language: Simpl.⁴

The Simpl Language

The Simpl language is defined as follows. The two basic type definitions are identifiers and labels—which are used to identify expressions in the analysis framework.

```
type Ident = String
type Label = Int
```

The most primitve expressions then are Atoms which, as in JavaScript, can be numbers, strings, regular expressions, booleans, identifiers—for variables—and the two values null and undefined.

Then we have expressions, which can represent all JavaScript's operators, and atomic expressions.

```
data Exp
-- * arithmetic expressions
= Add Exp Exp | Sub Exp Exp
| Mul Exp Exp | Div Exp Exp
| Mod Exp Exp | Neg Exp
```

⁴The implementation of the HJS to Simpl conversion algorithm can be found in the module Language.Javascript.Hjs2Simpl.

```
-- * boolean expressions
| And Exp Exp | Or
                     Exp Exp
Not Exp
            | Ter Exp Exp Exp
-- * comparison operators
| Eq Exp Exp | Neq Exp Exp
| SEq Exp Exp | SNeq Exp Exp
      Exp Exp | Gte Exp Exp
| Lt | Exp Exp | Lte | Exp Exp
-- * bitwise operators
| BAnd Exp Exp | BOr Exp Exp
BXor Exp Exp | BNot Exp
| BLs Exp Exp | BRs0 Exp Exp
BRs1 Exp Exp
-- * atomic expressions
Atom Atom
```

Expressions can be combined in "code", which represents expressions that can optionally have side-effects. Code instances are labelled.

Code instances can be further combined in statements, which contains all controlflow effecting expressions (aside from function calls).

Finally, statements—together with function declarations—compose a program.

```
data Decl = Decl Ident [Ident] Label Stmt Label
data Program = Program [Decl] [Stmt]
```

Note that because of this architecture, many JavaScript structures such as e.g. nested function declarations, functions-as-values, objects, etc... are unsupported.

Analysing Simpl